

Hidden Surface Elimination: BSP trees

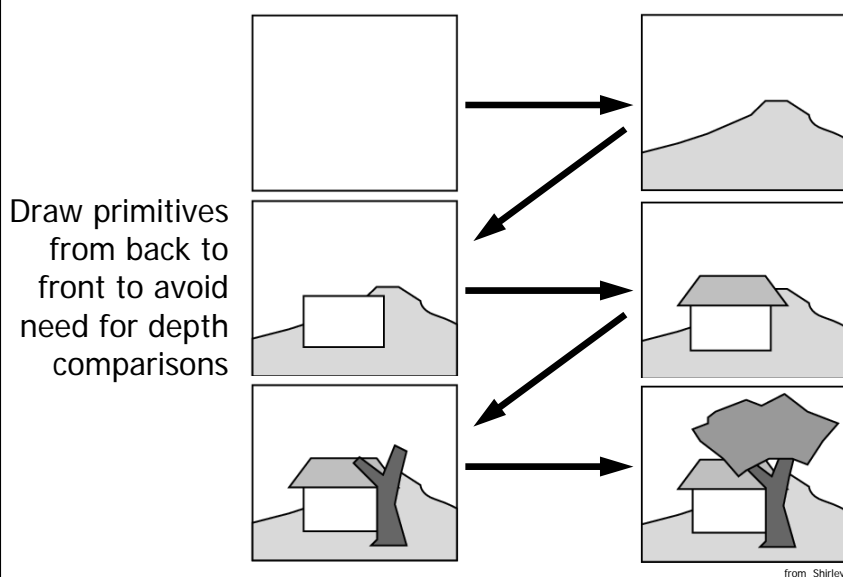
Outline

- Binary space partition (BSP) trees
 - Polygon-aligned

BSP Trees

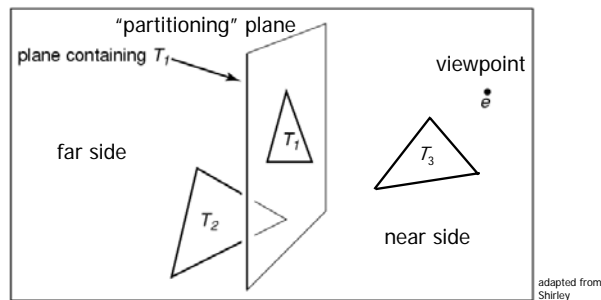
- Basic idea: Preprocess geometric primitives in scene to build a spatial data structure such that tests from **any viewpoint** can be easily calculated later
- Examples of tests
 - **Visibility** for painter's algorithms
 - **Intersection testing** for ray tracing
- Generalization of binary search trees (1-D) to higher dimensions

Painter's algorithm



BSP trees: Key property

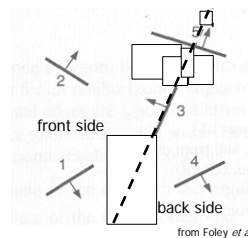
- “Spatial sorting” keeps track of which side of lines/planes primitives are on
 - Objects on the same side as the viewer can be **drawn on top** of objects on the opposite side
 - Objects on one side **cannot intersect** objects on the other side



- “Polygon-aligned” means partitioning plane is always coplanar with a scene polygon, as opposed to arbitrarily positioned

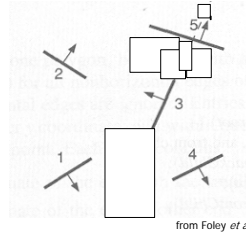
Building 2-D “line-aligned” BSP trees

- Pick oriented line segment (i.e., has a normal) from list as the root
- Rest of lines partitioned according to which side they are on
 - “Partitioning” line placed at root of subtree
 - Sets of lines on “front” side and “back” side correspond to left & right subtrees, respectively
- Recurse on each child

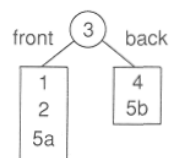
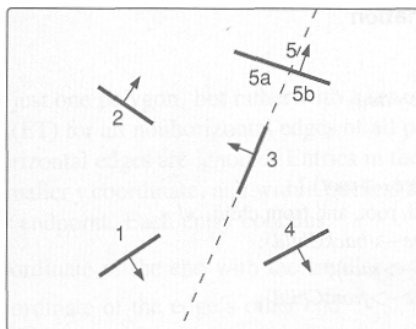


Building 2-D BSP trees: Issues

- How to pick line with which to partition
- What to do with lines that cross partitioning line
 - Split them (standard)
 - Or: Put a copy on each side of boundary
- When to stop recursing
 - n or fewer primitives per leaf ($n = 1$ is standard)
 - Or: Threshold on recursion depth

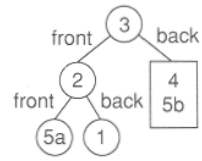
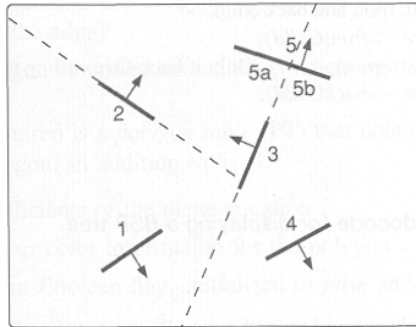


2-D BSP tree: Building example



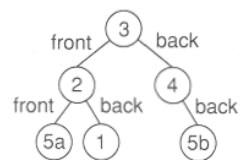
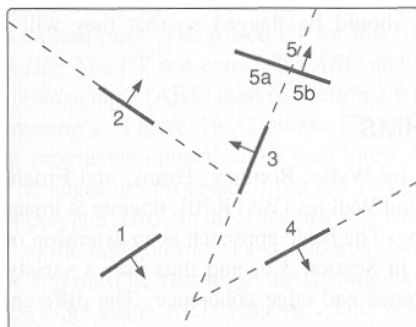
from Foley et al.

2-D BSP tree: Building example



from Foley *et al.*

2-D BSP tree: Building example



from Foley *et al.*

Building 2-D BSP trees: Pseudocode

adapted from Foley *et al.*

```
BSP_tree *BSP_makeTree(line *lineList)
{
    line root, l, backPart, frontPart, *backList, *frontList;

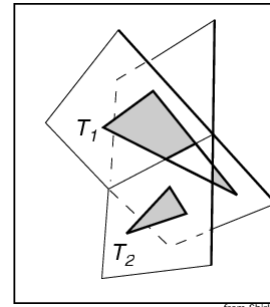
    if (lineList == NULL)
        return NULL;
    else {
        root = BSP_selectAndRemoveLine(lineList);
        backList = frontList = NULL;
        for (each line l in lineList) {
            if (l in front of root)
                BSP_addToList(l, frontList);
            else if (l in back of root)
                BSP_addToList(l, backList);
            else {
                BSP_splitLine(l, root, frontPart, backPart);
                BSP_addToList(frontPart, frontList);
                BSP_addToList(backPart, backList);
            }
        }
        return BSP_mergeTree(BSP_makeTree(frontList), root, BSP_makeTree(backList));
    }
}
```

Picking partitioning lines: Criteria

- **Painter's algorithm**
 - Every object must be drawn → Entire tree is traversed
 - So overall tree size should be as small as possible
 - Minimize splitting
- **Ray tracing** (later in course...)
 - Several paths from root to leaves traversed looking for intersections
 - So tree depth more important than overall size
 - Balance primitives on each side

Partitioning lines: "Least-crossed" heuristic

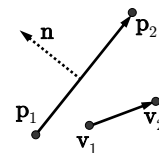
- For painter's algorithm, we want to choose partitioning lines that minimize the number of splits
- A particular line can be tested to see how many lines cross it and therefore would have to be split if it were the partitioning line
- Good procedure in practice:
 1. Randomly select a small number of candidate partitioning lines (e.g., 5-10 out of 1,000)
 2. Calculate number of lines that cross each candidate
 3. Use candidate with least crossings as next partition



from Shirley
Crossings for two planes

Building 2-D BSP trees: Details

- How to parametrize partitioning line \mathbf{l} from line segment $\mathbf{p}_1\mathbf{p}_2$?
 - Homogeneous line form $\mathbf{l} = (a, b, c)^T$ for segment is same as line it's on
 - We can obtain \mathbf{l} from $\mathbf{p}_1, \mathbf{p}_2$ via $\mathbf{l} = \mathbf{p}_1 \times \mathbf{p}_2$
 - Normal vector of line \mathbf{l} is $\mathbf{n} = (a, b)^T$
- How to decide which side of partitioning line \mathbf{l} a line segment $\mathbf{v}_1\mathbf{v}_2$ is on?
 - If both points $\mathbf{v}_1, \mathbf{v}_2$ are on the front side of \mathbf{l} , line $\mathbf{v}_1\mathbf{v}_2$ is on front side
 - If both $\mathbf{v}_1, \mathbf{v}_2$ are on the back side, $\mathbf{v}_1\mathbf{v}_2$ is on back side
 - If $\mathbf{v}_1, \mathbf{v}_2$ are on different sides, $\mathbf{v}_1\mathbf{v}_2$ crosses partitioning line \mathbf{l}

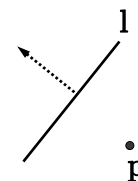


Building 2-D BSP trees: Details

- How to split crossing lines?
 - Find intersection point \mathbf{x} of $\mathbf{v}_1\mathbf{v}_2$ with \mathbf{l}
 - Let \mathbf{l}' be homogeneous form of line defined by $\mathbf{v}_1\mathbf{v}_2$
 - By definition, we want a point \mathbf{x} that is on both lines \mathbf{l} and \mathbf{l}' . This would imply that $\mathbf{l} \cdot \mathbf{x} = \mathbf{l}' \cdot \mathbf{x} = 0$
 - Just looking at these as vectors, a dot product of 0 means that \mathbf{x} is orthogonal to both \mathbf{l} and \mathbf{l}'
 - Because **cross product** is orthogonal to both multiplicands, $\mathbf{x} = \mathbf{l} \times \mathbf{l}'$ satisfies this requirement and thus defines the point of intersection
 - Given intersection \mathbf{x} :
 - If \mathbf{v}_1 is on front side of \mathbf{l} : Output $\mathbf{v}_1 \mathbf{x}$ as front part and \mathbf{xv}_2 as back part
 - If \mathbf{v}_2 is on front side of \mathbf{l} : Output \mathbf{xv}_2 as front part and $\mathbf{v}_1 \mathbf{x}$ as back part

Painter's algorithm: Tree traversal

- Want farthest-to-nearest ordering of primitives for painter's algorithm
 - If view location is on **front side** of a partitioning line:
 - Lines on back side are farther
 - Lines on front side are nearer
 - If view location is on **back side** of a partitioning line:
 - Lines on front side are farther
 - Lines on back side are nearer
- Which side of a partitioning line \mathbf{l} is a point \mathbf{p} on?
 - Assuming \mathbf{l}, \mathbf{p} in homogeneous form, use homogeneous line test $F(\mathbf{p}) = \mathbf{l} \cdot \mathbf{p}$
 - $F > 0 \Rightarrow \mathbf{p}$ is on **back side** of \mathbf{l}
 - $F < 0 \Rightarrow \mathbf{p}$ is on **front side** of \mathbf{l}
 - $F = 0 \Rightarrow \mathbf{p}$ is on line \mathbf{l} (arbitrarily treat as back side)

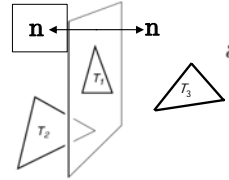


Painter's algorithm 2-D BSP tree traversal: Pseudocode

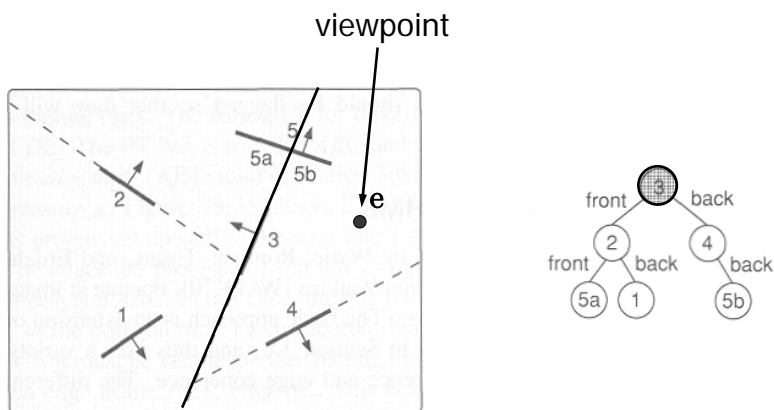
```

void BSP_displayTree(BSP_tree *tree, point viewLocation)
{
    if (tree != NULL) {
        if (viewLocation is in front of tree->root) {
            // Display back child, root, then front child
            BSP_displayTree(tree->backChild, viewLocation);
            displayLine(tree->root);
            BSP_displayTree(tree->frontChild, viewLocation);
        }
        else {
            // Display front child, root, then back child
            BSP_displayTree(tree->frontChild, viewLocation);
            displayLine(tree->root); // back-facing line-can cull by skipping
            BSP_displayTree(tree->backChild, viewLocation);
        }
    }
}

```



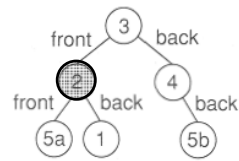
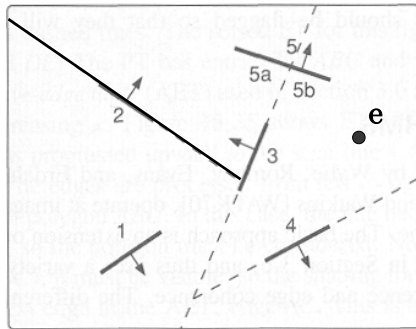
Painter's algorithm: Example BSP tree traversal



from Foley *et al.*

Behind root (node 3), so display everything in front of (left subtree = nodes 1, 2, 5a), then root (node 3), then everything behind (right subtree = nodes 4 and 5b)

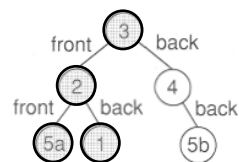
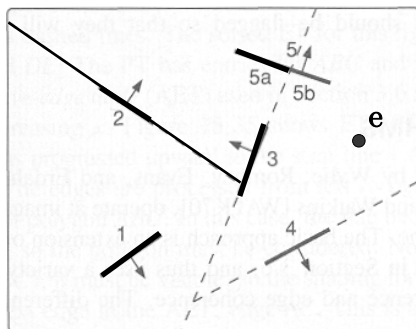
Painter's algorithm: Example BSP tree traversal



from Foley *et al.*

In front of root (node 2), so display everything behind (right subtree = node 1), then root (node 2), then everything in front of (left subtree = node 5a)

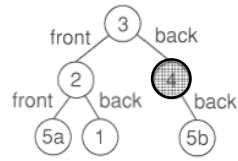
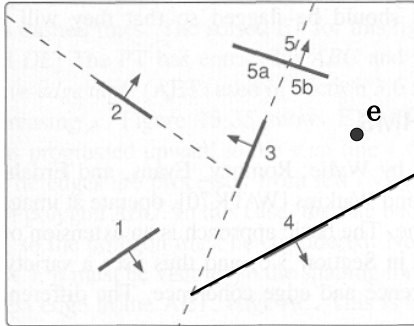
Painter's algorithm: Example BSP tree traversal



from Foley *et al.*

In front of root (node 2), so display everything behind (right subtree = node 1), then root (node 2), then everything in front of (left subtree = node 5a)

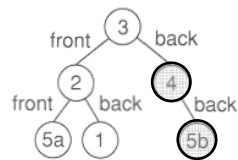
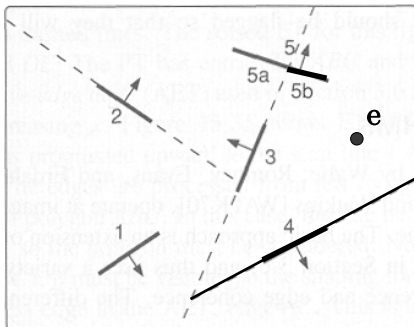
Painter's algorithm: Example BSP tree traversal



from Foley *et al.*

Behind root (node 4), so display everything in front of (left subtree = NULL), then root (node 4), then everything behind (right subtree = node 5b)

Painter's algorithm: Example BSP tree traversal

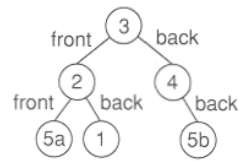
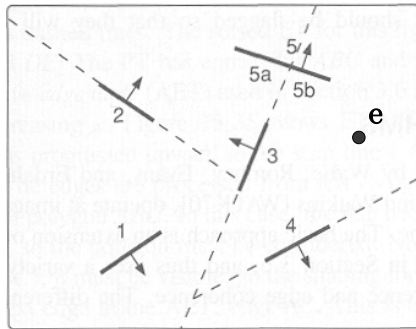


from Foley *et al.*

Behind root (node 4), so display everything in front of (left subtree = NULL), then root (node 4), then everything behind (right subtree = node 5b)

Painter's algorithm: Example BSP tree traversal

Final order: 1, 2, 5a, 3, 4, 5b

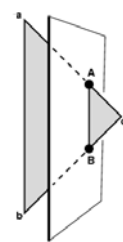


from Foley *et al.*

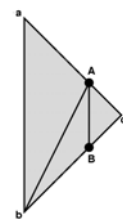
Every node is visited from back-to-front, so this is an $O(n)$ operation (n is the number of primitives **after** splitting)

3-D BSP Trees

- Analog of 2-D method, but now we are talking about 3-D polygon primitives and partitioning planes
- What's different about this vs. lines?
 - Must parametrize plane from polygon
 - Point-plane "sidedness" test is analogous to point-line test
 - Just use homogeneous form of plane equation
 - Line (each edge of polygon)-plane intersection instead of line-line intersection
 - Polygon splitting instead of line splitting



Triangle crossing partitioning plane



Triangle splitting

from Shirley

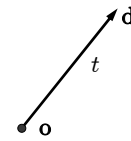
3-D BSP Trees: Details

- Parametrize plane from polygon
 - Cross product of edges gets plane normal $\mathbf{n} = (a, b, c)^T$
 - Solve for d from single point on plane $(x, y, z)^T$ via plane equation $ax + by + cz + d = 0$

- Line-plane intersection

- Unfortunately, there is no simple homogeneous form for lines in 3-D like in 2-D
- Instead:

- Express point \mathbf{p} on a ray as some distance t along direction \mathbf{d} from origin \mathbf{o} : $\mathbf{p} = \mathbf{o} + t\mathbf{d}$
- Use plane equation $\mathbf{n} \cdot \mathbf{x} + d = 0$, substitute $\mathbf{o} + t\mathbf{d}$ for \mathbf{x} , and solve for t
- Then plug t back into $\mathbf{p} = \mathbf{o} + t\mathbf{d}$ to get \mathbf{p}



BSP Trees: Notes

- Works best for static scenes
 - Moving primitives can cross partitioning lines
 - Dynamic adjustment of tree possible, but slows things down